

Tutorial 2

Exercise 1. For each $n \in \mathbb{N}$, define $S_n = 1 \oplus 2 \oplus \cdots \oplus n$.

(i) Find the values of S_{33} and S_{34} .

(ii) Consider the n -pile nim game with position $(1, 2, \dots, n)$.

(1) Find the values of n such that the position $(1, 2, \dots, n)$ is a P-position.

(2) Find all winning moves from the position $(1, 2, \dots, n)$ for $n = 33$.

(3) Find all winning moves from the position $(1, 2, \dots, n)$ for $n = 34$.

Solution: (i) By simple calculation, we have

n	1	2	3	4	5	6	7	8	9	10	11	12...
S_n	1	3	0	4	1	7	0	8	1	11	1	12...

It is easy to prove by induction that

$$S_n = \begin{cases} n & \text{if } n \equiv 0 \pmod{4} \\ 1 & \text{if } n \equiv 1 \pmod{4} \\ n + 1 & \text{if } n \equiv 2 \pmod{4} \\ 0 & \text{if } n \equiv 3 \pmod{4} \end{cases}.$$

Since $33 \equiv 1 \pmod{4}$ and $34 \equiv 2 \pmod{4}$, we have $S_{33} = 1$ and $S_{34} = 35$.

(ii). (1) The set of P-positions is

$$\{(1, 2, \dots, 4k + 3) : k = 0, 1, \dots\}.$$

(2) Since

$$\begin{array}{r}
 (0, 0, 0, 0, 0, 1)_2 \\
 (0, 0, 0, 0, 1, 0)_2 \\
 \vdots \\
 (1, 0, 0, 0, 0, 0)_2 \\
 (1, 0, 0, 0, 0, 1)_2 \\
 \hline
 (0, 0, 0, 0, 0, 1)_2 = 1
 \end{array}
 ,$$

we have all winning moves are: $(1, 2, \dots, k-1, k, k+1, \dots, 33) \rightarrow (1, 2, \dots, k-1, k-1, k+1)$ for all odd number k .

(3) Since

$$\begin{array}{r}
 (0, 0, 0, 0, 0, 1)_2 \\
 (0, 0, 0, 0, 1, 0)_2 \\
 \vdots \\
 (1, 0, 0, 0, 0, 0)_2 \\
 (1, 0, 0, 0, 0, 1)_2 \\
 (1, 0, 0, 0, 1, 0)_2 \\
 \hline
 (1, 0, 0, 0, 1, 1)_2 = 35
 \end{array}
 ,$$

we have all winning moves are: $(1, 2, \dots, 32, 33, 34) \rightarrow (1, 2, \dots, 3, 33, 34)$
or $(1, 2, \dots, 33, 34) \rightarrow (1, 2, \dots, 32, 2, 34)$, or $(1, 2, \dots, 33, 34) \rightarrow (1, 2, \dots, 33, 1)$.

Sprague-Grundy function.

Definition 1. Let X be the set of all possible positions of a combinatorial game. The S-G function is a map $g : X \rightarrow \mathbb{N}$ defined by

(i) $g(x) := 0$ if x is a terminal position.

(ii) $g(y) = \min\{k \geq 0, k \notin \{g(x) : x \text{ is a follower of } y\}\}$.

The S-G function is a useful tool designed to find the P-positions of a game, as we have the following proposition.

Proposition 2. *Let \mathcal{P} denote the set of P-positions and let g be the S-G function of a game. Then we have*

$$\mathcal{P} = \{x \in X : g(x) = 0\}.$$

Exercise 2. *Consider the subtraction game with $S = \{1, 3, 6\}$.*

- (i) *Find $g(6)$, $g(13)$ and $g(50)$.*
- (ii) *Find all winning moves from the position that there are 50 chips.*
- (iii) *Find the set of P-positions and give a proof.*

Solution (i) Note that the only terminal position is 0. By backwards induction, we have

x	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	...
$g(x)$	0	1	0	1	0	1	2	3	2	0	1	0	1	0	1	2	3	2	...

It is easy to see that g is periodic with period 9. Indeed,

$$g(x) = \begin{cases} 0 & \text{if } x \equiv 0, 2 \text{ or } 4 \pmod{9} \\ 1 & \text{if } x \equiv 1, 3 \text{ or } 5 \pmod{9} \\ 2 & \text{if } x \equiv 6 \text{ or } 8 \pmod{9} \\ 3 & \text{if } x \equiv 7 \pmod{9} \end{cases}$$

Hence we have $g(6) = 2$, $g(13) = 0$ and $g(50) = 1$.

- (ii) All winning moves from position 50 are removing 1 or 3 chips.

(iii) We claim that the set of P-positions is given by

$$\mathcal{P} = \{k \in \mathbb{N} : k \equiv 0, 2 \text{ or } 4 \pmod{9}\}.$$

Proof of the claim: (1). The only terminal position $k = 0$ is in \mathcal{P} . (2). If $k \equiv 0, 2$ or $4 \pmod{9}$, then we have $k - 1 \equiv 8, 1$ or $3 \pmod{9}$, $k - 3 \equiv 6, 8$ or $1 \pmod{9}$ and $k - 6 \equiv 3, 5$ or $7 \pmod{9}$. Hence any position in \mathcal{P} can only be removed to a position outside \mathcal{P} . (3). If $k \not\equiv 2 \pmod{9}$, $k \not\equiv 4 \pmod{9}$ and $k \not\equiv 6 \pmod{9}$, it is easy to see that at least one of $k - 1$, $k - 3$ and $k - 6$ is in \mathcal{P} . By the characterization of P-positions, we finish the proof.